



## SSD Architectures and Experiences

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- Flash Overview
- Performance Characterization Methodology
- Technology Overview

## Virident Corporate Overview



- Mission: Unconditional Performance Technology
  - Experts at system-level NAND management for performance and endurance
  - Blue-chip investors: Sequoia Capital, Globespan Capital, Artiman, Intel Capital, Cisco Ventures
  - Over 100 technical staff in Milpitas, CA and Bangalore, India
- Product: FlashMAX storage class memory
  - Unconditional performance: Highest IOs/G, Lowest latency (µs), Sustained over time
  - Enterprise class reliability: Guaranteed writes, ECC/RAID on board, 7+1 distribution
  - Highest Density: PCIe, half width, half height
  - Delivers 2-5X <u>application-level</u> sustained performance<sup>1</sup>
- Customers: High performance datacenter storage
  - Computer and Storage OEMs, Web Portals, Cloud Providers, Fortune 1000 Enterprise
  - U.S., Japan, APAC

<sup>1</sup>Based on 4K Random 70% Read, 30% Write – the typical TPC-C DB workload model,

# **Traditional Storage Challenges**





#### Performance gap between HDD and DDR

## Flash Storage's Promise

### Filling the Performance Gap





#### **Performance of devices continues to increase**





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**Performance Characterization Methodology** 



- Level 1: Initial Performance (SNIA FOB)
- Level 2: Baseline performance (SNIA Transition)
- Level 3: Sustained performance, IO-QoS metrics (SNIA Steady State)
- Level 4: Application performance and real-world workloads





## **Goal: Highest, Sustained Performance**

Across Wide Range of Applications





Example performance for sample of solid state devices

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What Makes Delivering "IO QoS" Difficult?



### NAND Flash characteristics:

- •Asymmetry
  - Difference between read and write granularity
  - Operation latencies
- Lack of update-in-place capability
  - Needs constant balancing between user traffic and flash management
- •Endurance, reliability
  - Worse with increasing capacity (or span)
- Each can have an order-of-magnitude impact...

## An IO-IOS flash architecture



Low resource overhead

### Sustained Performance Mechanisms Scheduling and Rate-Controlled Garbage Collection





## Low-level benchmark: IOPS





Test command: "run\_all.sh 1 1800 0"

Innosim rev 78: http://bazaar.launchpad.net/~mdcallag/mysql-patch/mytools/files/head:/bench/innosim

CPU: 2 x Intel X5690 (6 cores @3.43GHz, HT on)

OS: CentOS 6, 2.6.32-220.el6

File system: XFS with 4KB sector size (mkfs.xfs -s size=4096). Full drive capacity pre-filled with data before creating file system.

## Low level benchmark: latencies





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## Application Benchmark: MySQL



Very Steady Performance, Scales with Threads







- Flash needs guaranteed high performance
- Requirements based on SNIA Steady-State and Application Benchmarks
- Architecture built to plug-n-play performance in worst case today
- Future product direction aligned with NVMe